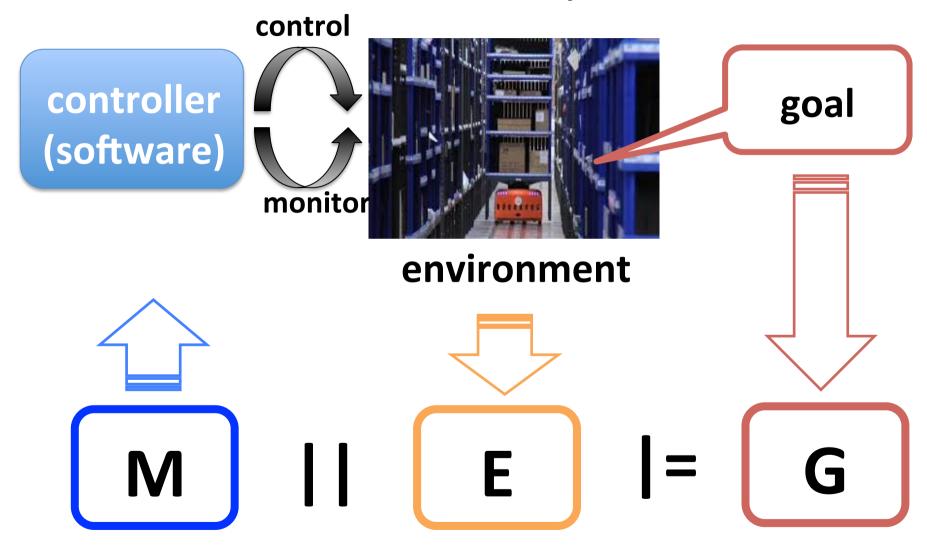
# Assured Graceful Degradation with Discrete Controller Synthesis

Kenji Tei National Institute of Informatics

Joint work with

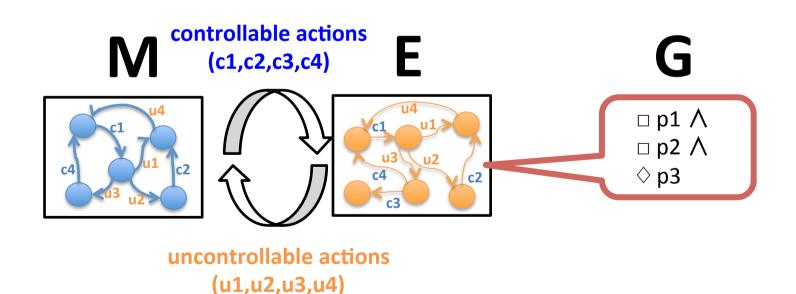
Kazuya Aizawa, Waseda University Nicolas D'Ippolito, Universidad de Buenos Aires

## Assurance at development time

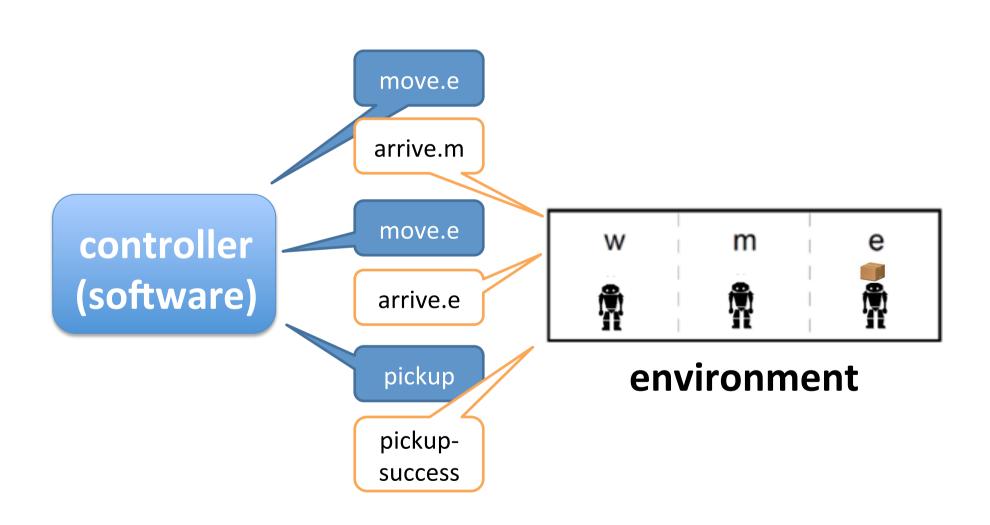


## Modeling approach: LTS and FLTL

- M and E: labeled transition system (LTS)
- G: fluent linear temporal logic (FLTL)



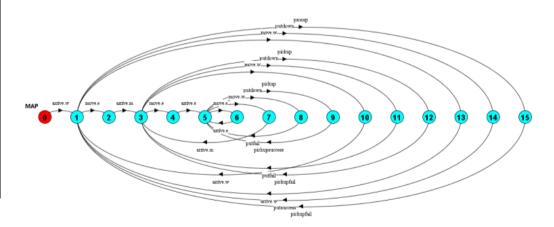
## **Example: Automated Warehouse**



## Modeling an environment by LTS

 $| | E = (MAP | W_ROBOT).$ 

```
MAP=(arrive['w] -> MAP['w]),
MAP['w]=( move['e] -> arrive['m] -> MAP['m]
| move['w] -> arrive['w] -> MAP['w]
| putdown -> putsuccess -> MAP['w]
| pickup -> pickupfail -> MAP['w]),
MAP['m]=( move['e] -> arrive['e] -> MAP['e]
| move['w] -> arrive['w] -> MAP['w]
| putdown -> putfail -> MAP['m]
| pickup -> pickupfail -> MAP['m],
MAP['e]=( move['e] -> arrive['e] -> MAP['e]
| move['w] -> arrive['m] -> MAP['m]
| putdown -> putfail -> MAP['e]
| pickup -> pickupsuccess -> MAP['e]).
```



Modeling how the state of the env. is changed and how the env. will react

## Specifying Goals by FLTL

[]((AT['w] && X(move['e])) -> X(!arrive['w] W pickupsuccess))

[]((AT['e] && X(move['w])) -> X(!arrive['e] W putsuccess))

[](putdown->AT['w])

[]!(!<pickupsuccess,putsuccess> && putdown)

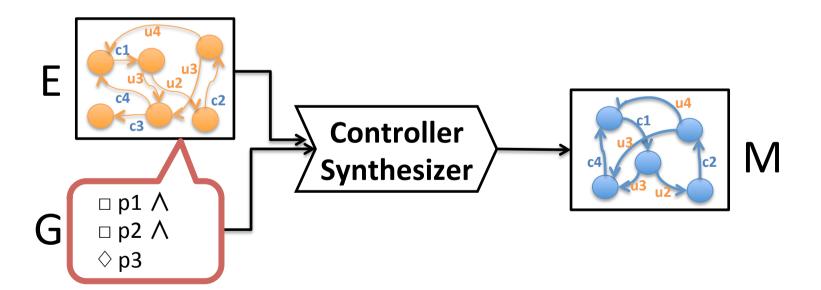
[](pickup->AT['e])

[]!(<pickupsuccess,putsuccess> && pickup)

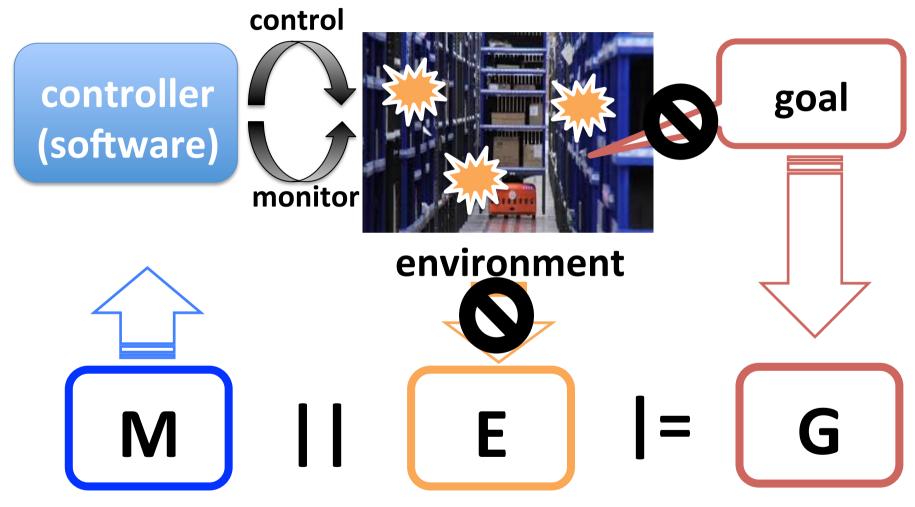
[](<ended,reset> -> (<pickupsuccess,{reset}> && <putsuccess,{reset}>))

## A way to generate M with assurance

- discrete controller synthesis [D'Ippolito, 2010] [D'Ippolito, 2011]
  - solve a control problem <E,G> to find an LTS M



## E may be invalid at runtime



System *may no longer work*, or *may continue*, but *without any assurances* 

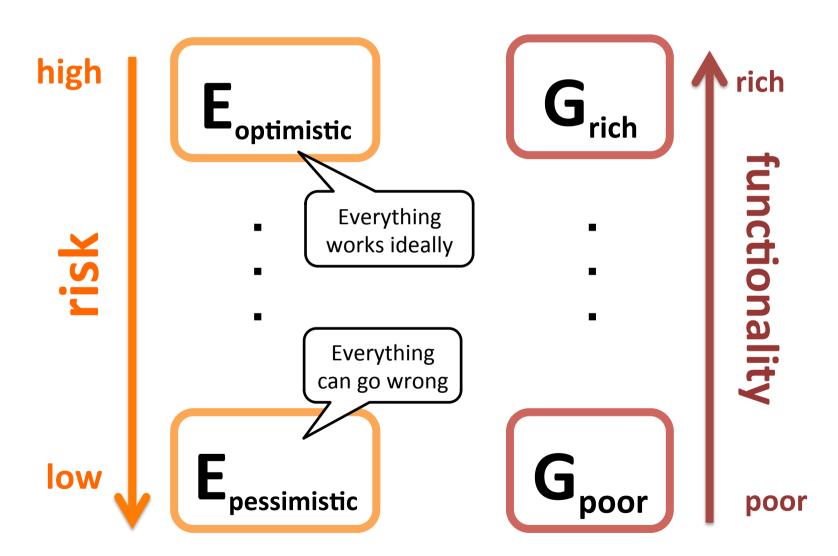
## **Assuming More Realistic**

```
[]((AT['w] && X(move['e])) -> X(!arrive['w] W pickupsuccess))
MAP['w]=( move['e] -> arrive['m] -> MAP['m]
                                                                                     []((AT['e] && X(move['w])) -> X(!arrive['e] W putsuccess))
 | move['w] -> arrive['w] -> MAP['w]
                                                                               [](putdown->AT['w])
                                                                                                     []!(!<pickupsuccess,putsuccess> && putdown)
 | putdown -> putsuccess -> MAP['w]
 | pickup -> pickupfail -> MAP['w]),
                                                                                [](pickup->AT['e])
                                                                                                      []!(<pickupsuccess,putsuccess> && pickup)
                                                                                [](<ended,reset> -> (<pickupsuccess,{reset}> && <putsuccess,{reset}>))
MAP['w]=(move['e] \rightarrow (arrive['m] \rightarrow MAP['m])
                                                                                     []((AT['e] && X(move['w])) -> X(!arrive['e] W putsuccess))
 | (arrive['w] -> MAP['w])
| move['w] -> arrive['w] -> MAP['w]
                                                                               [](putdown->AT['w])
                                                                                                     []!(!<pickupsuccess,putsuccess> && putdown)
 | putdown -> putsuccess -> MAP['w]
                                                                                [](pickup->AT['e])
                                                                                                      []!(<pickupsuccess,putsuccess> && pickup)
```

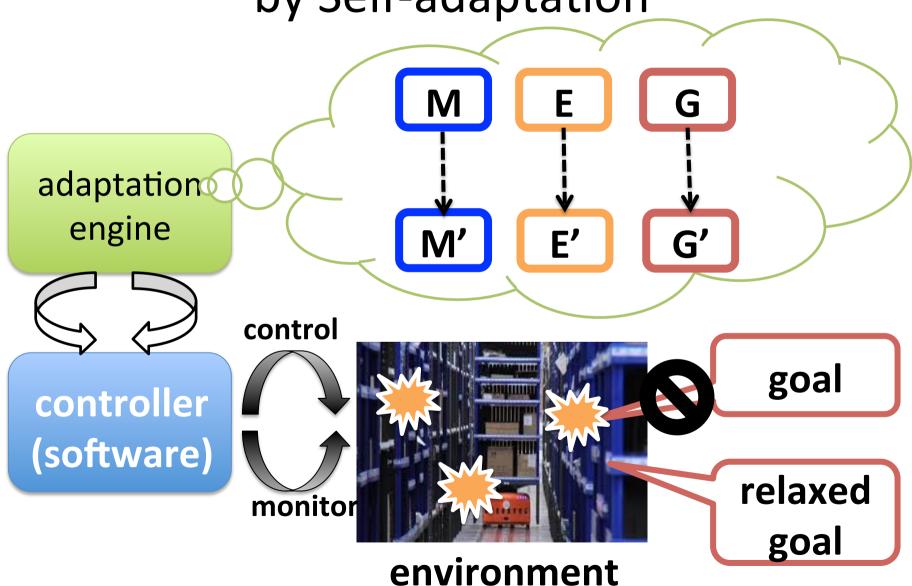
[](<ended,reset> -> (<pickupsuccess,{reset}> && <putsuccess,{reset}>))

| pickup -> pickupfail -> MAP['w]),

## How much should we assume?



Graceful Degradation by Self-adaptation



## Questions

How can the system be made to degrade gracefully with assurance?

How can the system determine how much it should degrade?

## Objective

We propose a framework for adaptation engine enabling graceful degradation

should not degrade the system too much

 should assure that the system after degradation satisfies a selected level of goals

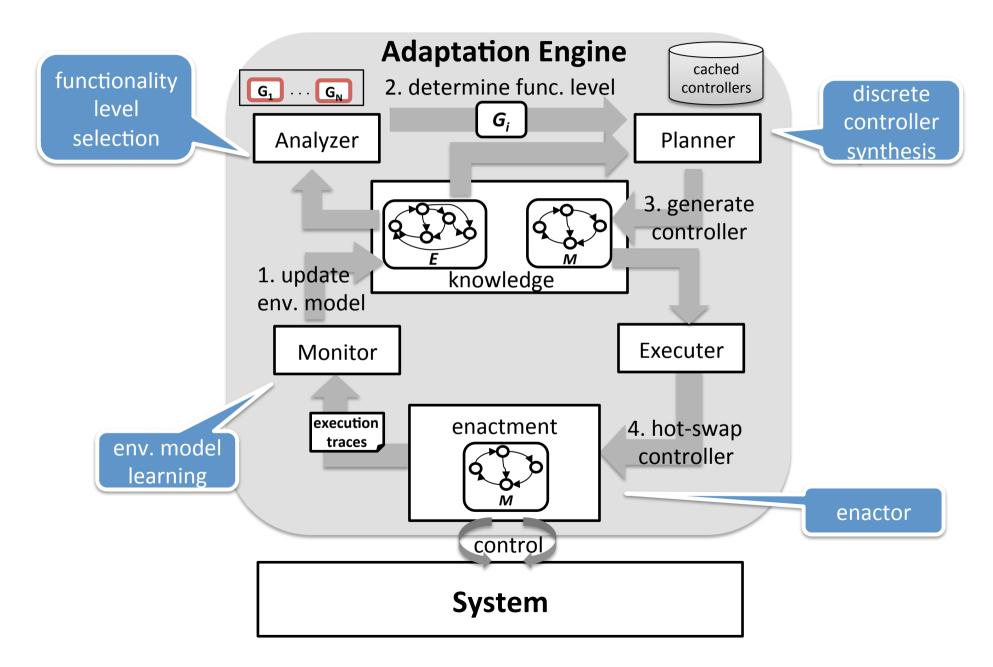
#### M' simulate M Performing degradation seamlessly

- should not stop or restart the system
- M' should simulate M

## Approach: Models@Run.time

- Revising environment model at runtime
  - to fit the environment
- Generating behavior specification with assurance at runtime
  - by using algorithmic techniques,
     in particular discrete controller synthesis
- Change behavior of the system in accordance with the generated model

### Overall architecture



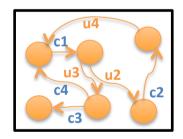
#### 0. Initialization

1. Specify levels of functionalities



2. Describe the initial environment model

 $E_{c}$ 



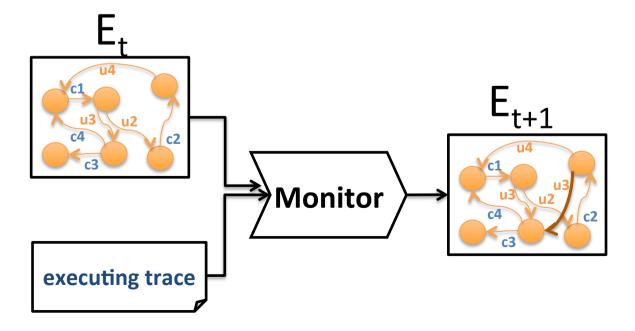
3. Select a func. level and construct the initial controller

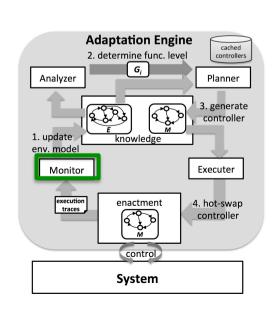




## 1. Monitor: Environment Model Updates

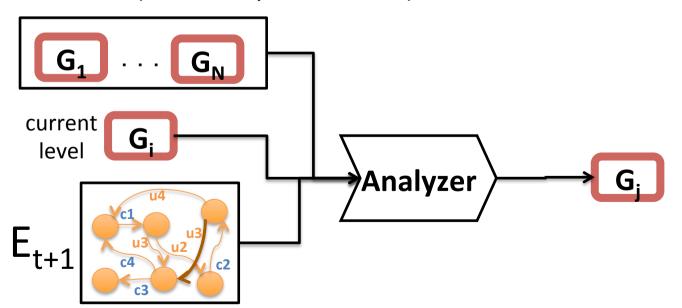
- update  $E_t$  to generate  $E_{t+1}$  so that  $E_{t+1}$  can explain execution traces of the system
  - find and add unmodeled uncontrollable transitions  $\Delta_{t+1}$ 
    - When a robot performed "move.w" action at "e", the environment will respond "arrive.m" or "arrive.e"
  - rule learning for environment model update [Sykes, 2013]

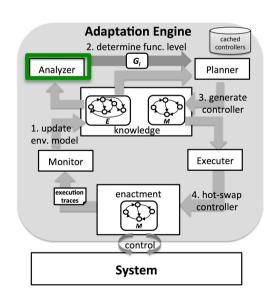




#### 2. Analysis: Functionality-level Selection

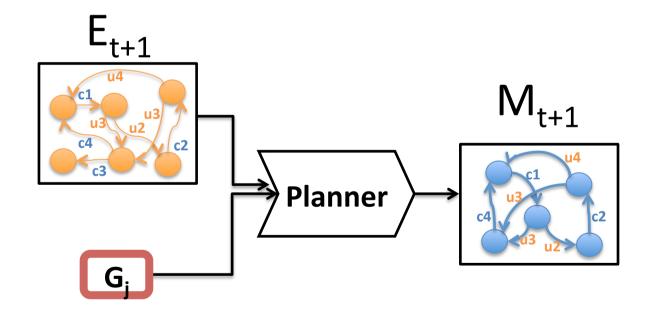
- Determine a functionality level  $G_j$  from  $\{G_1,...G_N\}$ 
  - $G_j$  can be satisfied in  $E_{t+1}$
  - The system can degrade to  $G_j$  without stopping or restarting itself
- Functionality level selection
  - (will be explained later)

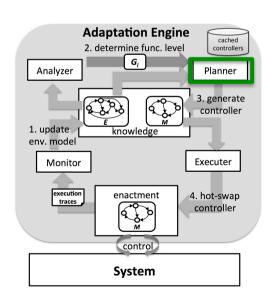




## 3. Plan: Discrete Controller Synthesis

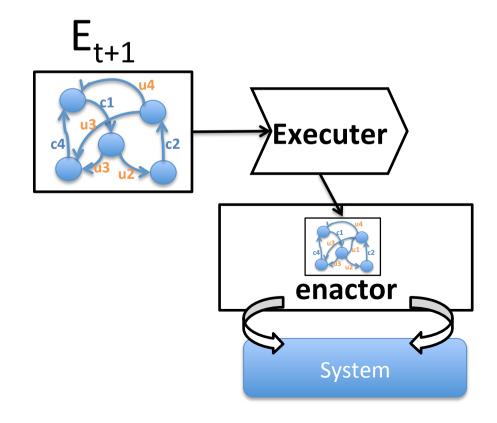
- Generate an LTS  $M_{t+1}$  guaranteeing satisfaction of  $G_j$  in  $E_{t+1}$
- Discrete controller synthesis [D'Ippolito, 2010] [D'Ippolito, 2011]
  - solve a control problem <E,G> to find an LTS M

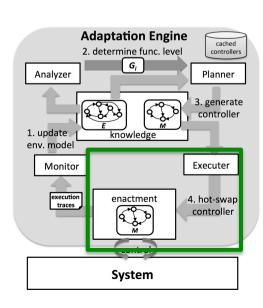




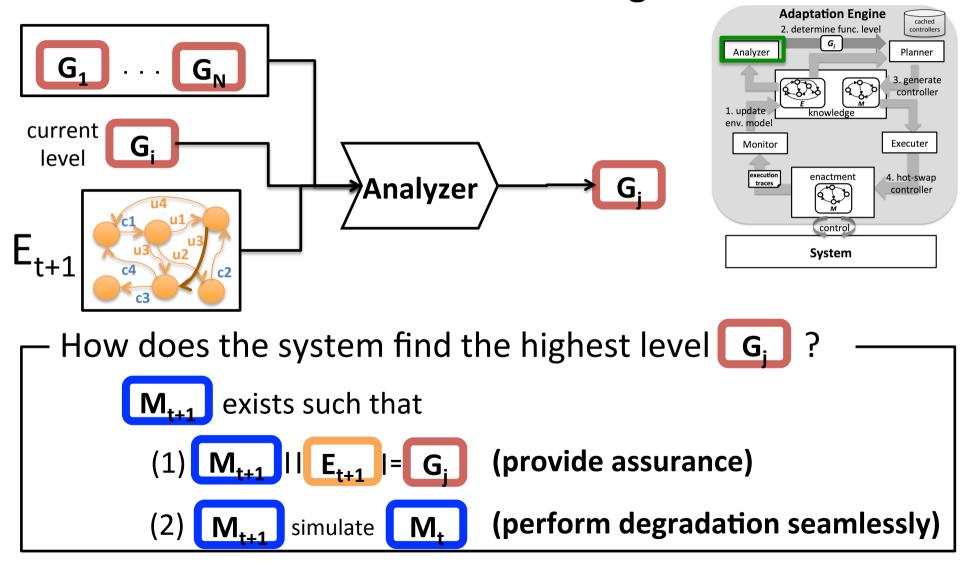
#### 4. Execute: Enactor

- Hot-swap controller model from  $M_t$  to  $M_{t+1}$ 
  - It can be done without stopping the system because  $M_{t+1}$  simulates  $M_t$
- Enactment framework [Braberman, 2013]
  - interpret LTS and orchestrate high-level operations provided by the system

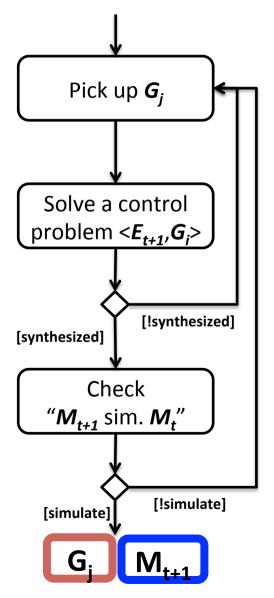




## Q: How can the system determine how much it should degrade?



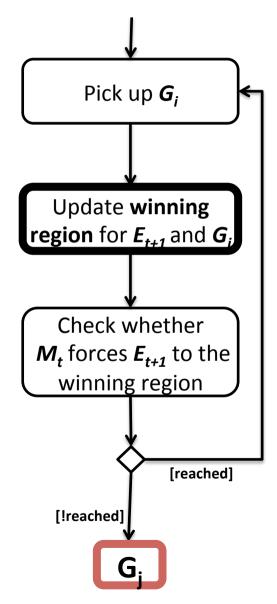
## Naive Strategy: Synthesize, then check



- Simple and straightforward
  - synthesized controller  $M_{t+1}$  can be used for the next controller

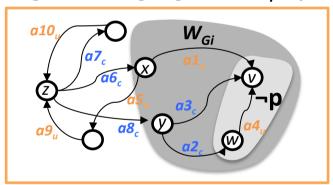
- Computationally inefficient
  - N control problems should be solved at worst

#### Advanced Strategy: Check without Synthesis



#### winning region W<sub>Gi,Et</sub>

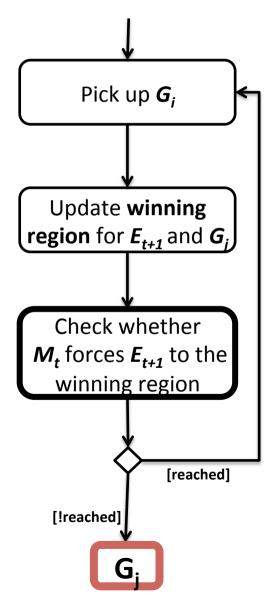
- the set of all states s such that no system forces  $E_t$  to satisfy  $G_i$  from s
- controller strategy should avoid the winning region e.g. a winning region for  $G_i = \Box p$

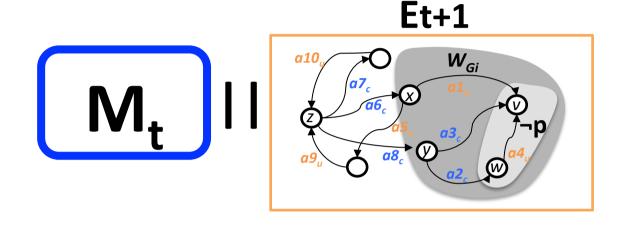


#### update winning region

- $W_{\textit{Gi,Et+1}}$  is obtained from  $W_{\textit{Gi,Et}}$  and  $\Delta_{\textit{t+1}}$
- determine states newly added in the region by checking updated part in env. model

#### Advanced Strategy: Check without Synthesis





- Does  $M_t | E_{t+1}$  reach to  $W_{Gi,Et+1}$ ?
  - if yes,  $M_{t+1}$  does not exist such that

$$M_{t+1}$$
 |  $E_{t+1}$  |  $E_{G_j}$ 
 $M_{t+1}$  simulate  $M_t$ 

## Case studies

automated warehouse

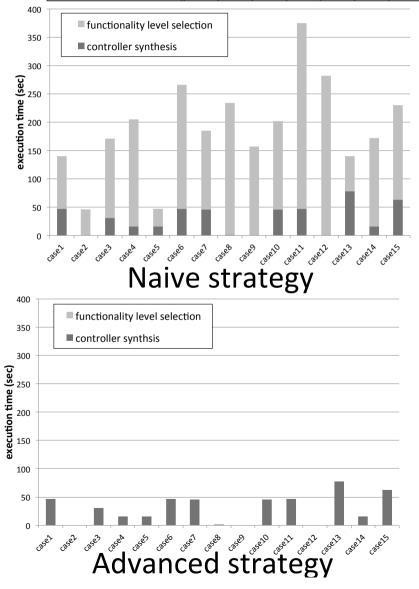


#### production cell



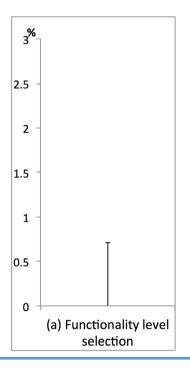
## **Automated Warehouse**

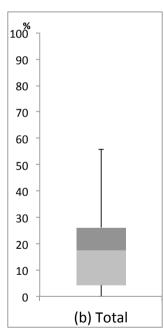
| Table 1: Case studies in the automated warehouse scenario |       |       |       |       |       |       |       |       |       |       |       |       |       |       |       |
|---|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|-------|
| Case  | 1     | 2     | 3     | 4     | 5     | 6     | 7     | 8     | 9     | 10    | 11    | 12    | 13    | 14    | 15    |
| before degradation  | $G_5$ | $G_5$ | $G_5$ | $G_5$ | $G_5$ | $G_4$ | $G_3$ | $G_4$ | $G_2$ | $G_5$ | $G_5$ | $G_5$ | $G_3$ | $G_2$ | $G_2$ |
| after degradation   | $G_5$ | $G_4$ | $G_3$ | $G_2$ | $G_4$ | $G_3$ | $G_1$ | $G_2$ | $G_1$ | $G_4$ | $G_2$ | $G_3$ | $G_3$ | $G_2$ | $G_1$ |
| # of levels checked                                       | 1     | 2     | 3     | 4     | 2     | 2     | 3     | 3     | 2     | 2     | 4     | 3     | 1     | 1     | 2     |



#### For func. selection

## For func. selection + controller synthesis





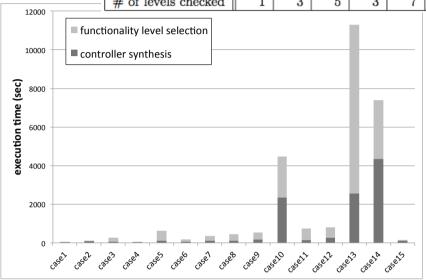
0.7% in the worst 0.00002% on average

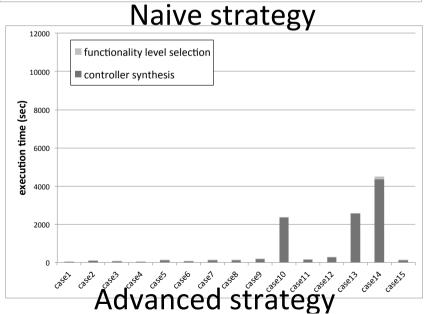
35.8% in the worst 13.6% on average

#### **Production Cell**

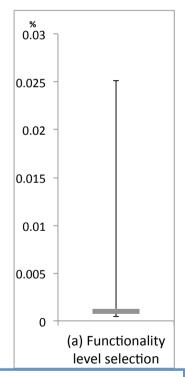
Table 2: Case studies in the production cell scenario

| Case                | 1        | 2        | 3        | 4        | 5        | 6        | 7        | 8        | 9     | 10    | 11    | 12    | 13       | 14    | 15       |
|---------------------|----------|----------|----------|----------|----------|----------|----------|----------|-------|-------|-------|-------|----------|-------|----------|
| before degradation  | $G_{12}$ | $G_{12}$ | $G_{12}$ | $G_{12}$ | $G_{11}$ | $G_{12}$ | $G_{10}$ | $G_{11}$ | $G_7$ | $G_5$ | $G_6$ | $G_4$ | $G_{11}$ | $G_6$ | $G_{11}$ |
| after degradation   | $G_{12}$ | $G_{10}$ | $G_8$    | $G_{10}$ | $G_5$    | $G_9$    | $G_7$    | $G_6$    | $G_4$ | $G_3$ | $G_2$ | $G_1$ | $G_6$    | $G_4$ | $G_{11}$ |
| # of levels checked | 1        | 3        | 5        | 3        | 7        | 4        | 4        | 6        | 4     | 3     | 5     | 4     | 6        | 3     | 1        |



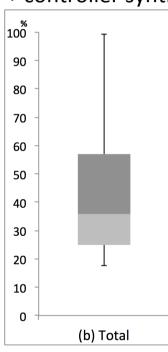


#### For func. selection



2.5% in the worst 0.176% on average

## For func. selection + controller synthesis



99.2% in the worst 44.9% on average

#### Conclusion

- How does the system cope with development time uncertainty?
  - How do we select appropriate level of functionality considering risks and functionality?
- We propose a framework enabling graceful degradation
  - revise environment model @ runtime
  - generate behavior specification with assurance @ runtime
  - change behavior of the system @ runtime
- We introduce two strategies to find the highest level of functionality that can be guarantee and to which the system can seamlessly degrade