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On the Comput. Complexity of the Dirichlet Problem for Poisson's Equation



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Recap: Complexity of Operators

- $f:[0;1] \rightarrow [0;1]$ polytime computable (\Rightarrow polyn. modulus of continuity)
- $f \rightarrow \int_0^1 f(t) dt$ (cmp. e.g. Valiant'79, Goldsmith&Ogihara&Rothe'98)
~~computable in exponential time;
 $\#P_1$ -"complete"~~ [Friedman&Ko'82]
 - $\int: f \rightarrow \int f: (x \rightarrow \int_0^x f(t) dt)$
~~computable in exponential time;~~
~~polytime-computable iff $\mathcal{FP} = \#P$~~
 - dsolve: $C[0;1] \times [-1;1] \ni f \rightarrow z: \dot{z}(t) = f(t, z), z(0) = 0.$
 - in general no computable solution $z(t)$
 - for $f \in C^1$ $PSPACE$ -"complete" [Kawamura'10, Kawamura et al]
 - for $f \in C^k$ CH -"hard"



Poisson's Equation

- [Pour-El&Richards'81, ..., Weihrauch&Zhong'02]
In-/computability of the *Wave Equation* (hyperbolic)
- Computability of some *nonlinear* PDEs:
[Yoshikawa'00;Weihrauch&Zhong]

PDE on connected

and open set Ω :

$$\Delta u = f \text{ on } \Omega, \quad u = g \text{ on } \partial\Omega$$

- electrostatic / gravitational potential of the charge/mass distribution f with boundary condition g
- 2nd order, linear, elliptic: homog. $(f,0)$ and inh. $(0,g)$
- 'fundamental' solutions $\ln |\underline{x}|$ (2D) and $1/|\underline{x}|$ (3D)
- 'explicit' Green's functions for various domains,
- solution formula on the complex unit disc; e.g. $g=0$

$$u(z) = -\frac{1}{2\pi} \cdot \int_{|w| \leq 1} \ln \frac{|w-z|}{|w \cdot z^* - 1|} \cdot f(w) dA(w)$$



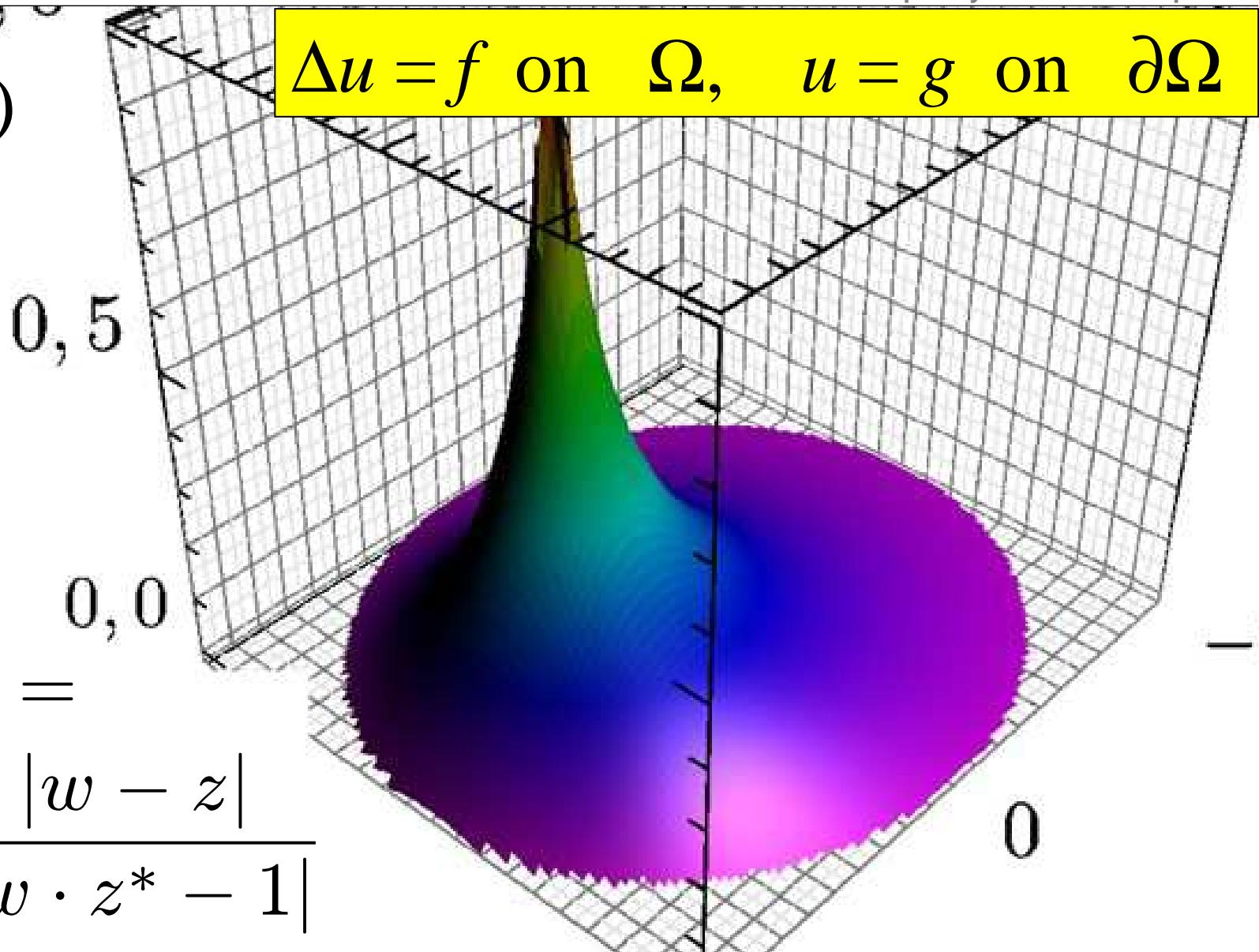
Green's Function in 2D

$G(1, w)$

$$\Delta u = f \text{ on } \Omega, \quad u = g \text{ on } \partial\Omega$$

$G(z, w) =$

$$\ln \frac{|w - z|}{|w \cdot z^* - 1|}$$



Results

$$\Delta u = f \text{ on } \Omega, \quad u = g \text{ on } \partial\Omega$$

Complexity of Poisson's Equation

Theorem: $B_d :=$ closed d -dim. Euklidean unit ball

- a) For every polytime $f:B_d \rightarrow \mathbb{R}$ and $g:\partial B_d \rightarrow \mathbb{R}$,
there exists a unique C^2 solution $u = \Phi(B_d, f, g)$
and u is computable in exponential time.
- b) If $\mathcal{FP} = \#\mathcal{P}$, then u is even polytime computable.
- c) There exists a polytime $f \in C^\infty$ such that
 $u = \Phi(B_d, f, 0)$ is polytime iff $\mathcal{FP} = \#\mathcal{P}$.
- d) For $d > 1$ there is a polytime $g \in C^\infty$ s.t.
 $\Phi(B_d, 0, g)$ is polytime iff $\mathcal{FP} = \#\mathcal{P}_1$

$$u(z) = -\frac{1}{2\pi} \cdot \int_{|w| \leq 1} \ln \frac{|w-z|}{|w \cdot z^* - 1|} \cdot f(w) dA(w)$$



Proof (Sketch)

- f, g polyn.modulus of continuity $\Rightarrow u = \Phi(f, g)$ is C^2 .
- Bound improper integral (singularity) for $w \rightarrow z$.
- Parameter integration over $|w-z| > 2^{-n}$ uniformly feasible in $\#\mathcal{P}$.
- Take polytime $h \in C^\infty[0;1]$ s.t. $\int h$ is " $\#\mathcal{P}$ -hard" and let $f(\underline{x}) := h'(|\underline{x}|) / |\underline{x}| \cdot \ln |\underline{x}|$ radially symmetric.

b) If $\mathcal{FP} = \#\mathcal{P}$, then u is even polytime computable.
there exists a unique C^2 solution $u = : \Phi(f, g)$
and u is computable in exponential time.

$$u(z) = -\frac{1}{2\pi} \cdot \int_{|w| \leq 1} \boxed{\ln \frac{|w-z|}{|w \cdot z^* - 1|}} \cdot f(w) dA(w)$$



Under further investigation...

Theorem: b+c) $f \rightarrow \Phi(B_d, f, 0)$ maps (smooth) polytime to polytime functions iff $\mathcal{FP} = \#\mathcal{P}$.

d) For $d > 1$ there is a polytime $g \in C^\infty$ s.t.

$\Phi(B_d, 0, g)$ is polytime iff $\mathcal{FP} = \#\mathcal{P}_1$

- Uniform: 2nd order polytime W-Reduction
- Close gap w.r.t. g between $\#\mathcal{P}_1$ and $\#\mathcal{P}$
- Consider domains other than balls B_d
- 2D → complexity of the Riemann mapping
- other (nonlinear) PDEs → Navier-Stokes!