Encoding and modeling for set compression

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Metadata about this work

- Started as broad attack on set compression
- Contribute one algorithm + smaller points
- Not particularly about data structures or time complexity, but setting goals/ideals and finding potentials

Events, sequences, sets

Event ("circumstance") sequence E_1, \ldots, E_n

$$\Pr(E_i|E_1 \cap \cdots \cap E_{i-1})$$

$$\sum_{i} -\log_2 \Pr(E_i|E_1 \cap \cdots \cap E_{i-1}) \text{ bits }$$

Encoder decoder share premises for what E_i mean in terms of specifying message (data)

- Seqence: a certain character is at position *i*
- Set: a certain element is in the set (note: no i)
- Set, alt.: a certain number of elements have a certain property

• . . .

Fields of application/ previous work

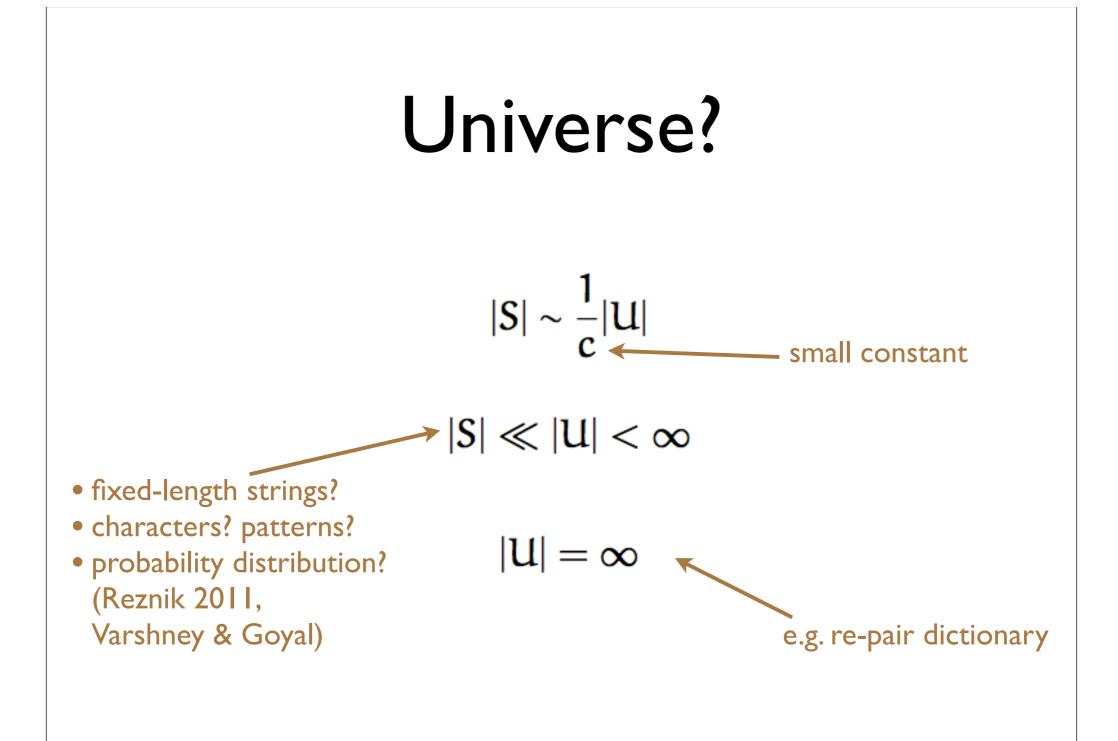
- Component (e.g. ψ)
- Inverted index
- Dictionary
- Data mining (measuring ratio), web graphs, ...

Universe?

Encoding a set (or many sets) S, elements drawn from universe U

 $S \in U$

|S| < |U|



Narrow focus, for now:

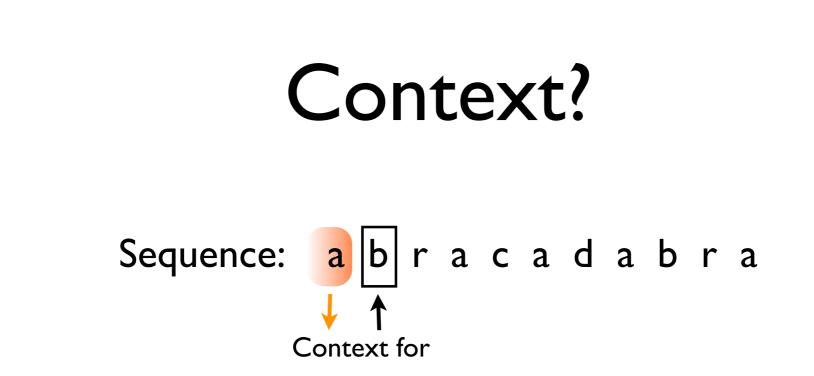
- U may be much larger than S
- Dependencies between elements
- Elements:
 - Integers $\in [0, |\mathbf{U}|)$
 - = bitstrings of length $\lceil \log_2 |U| \rceil$

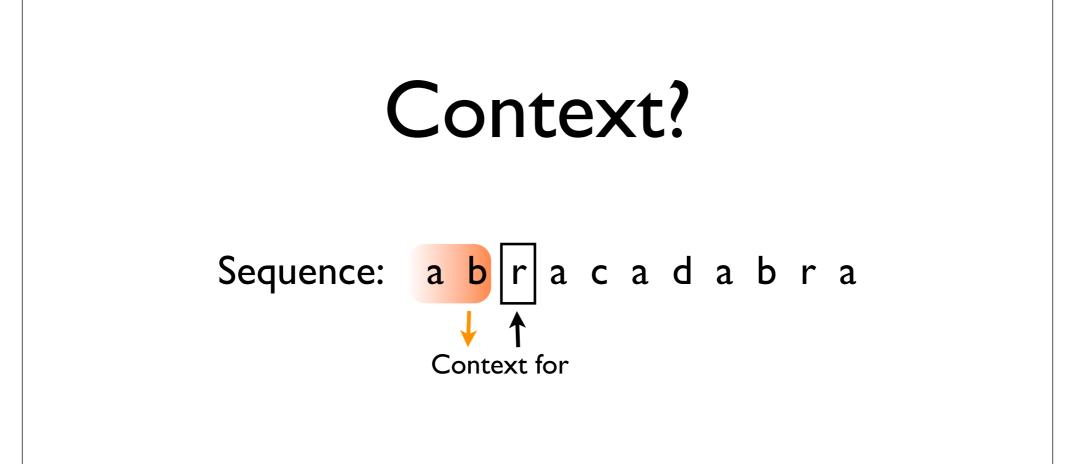
Solved?

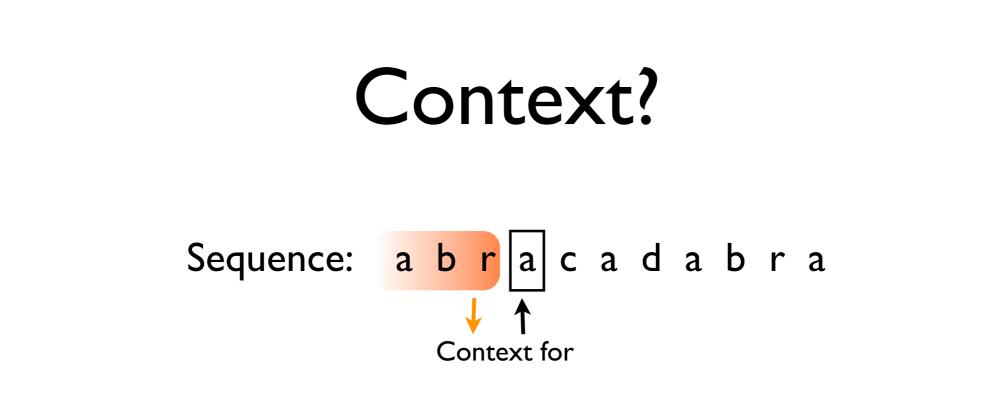
- Set: { 4, 9, 11, 14, 16, 17, 20, 21 }
- Gaps: { 4, 4, 1, 2, 1, 0, 2, 0 }
- Geometric distribution $Pr(gap size k) = (1-p)^{k-1}p, \ p = Pr(x \in S) = |S|/|U|$ Optimal code: Golomb (or arithmetic)

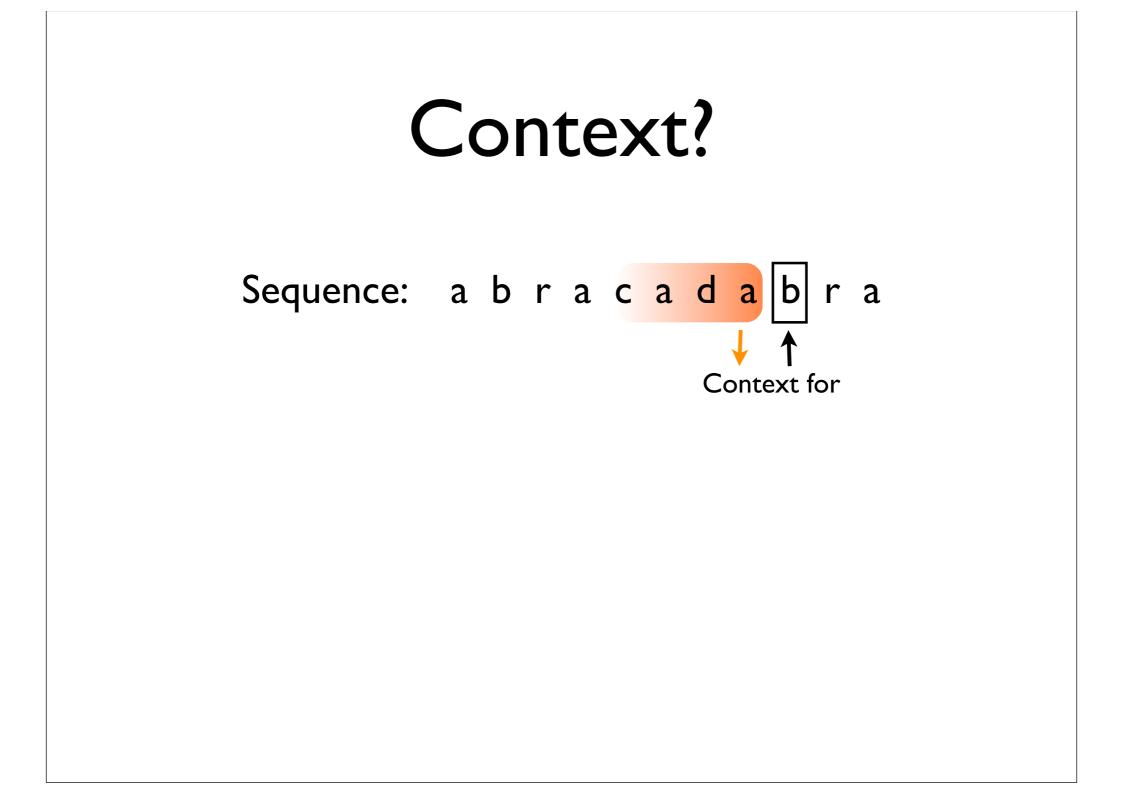
Known (?) method 2: yes/no code

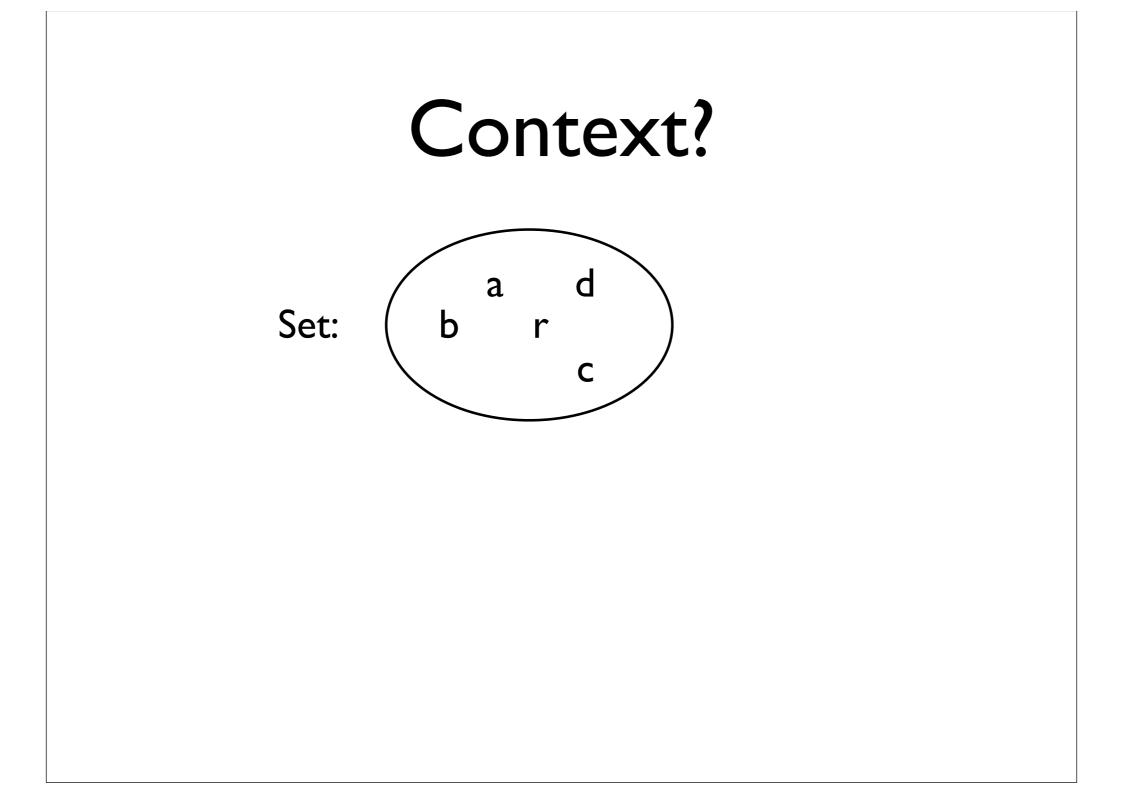
- Arithmetic code for binary source: for each element of U, encode whether in S
- Estimate $p_x = \Pr(x \in S)$
- Use probability ranges $[0, p_x), [p_x, 1]$

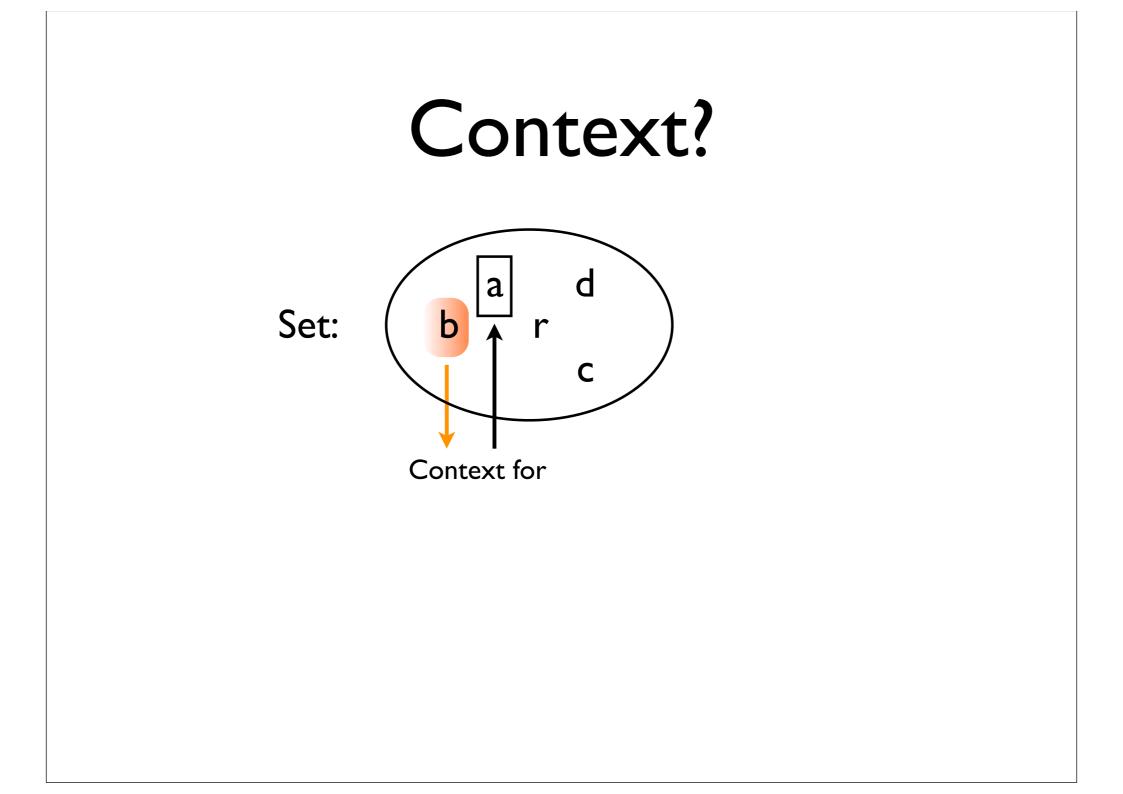


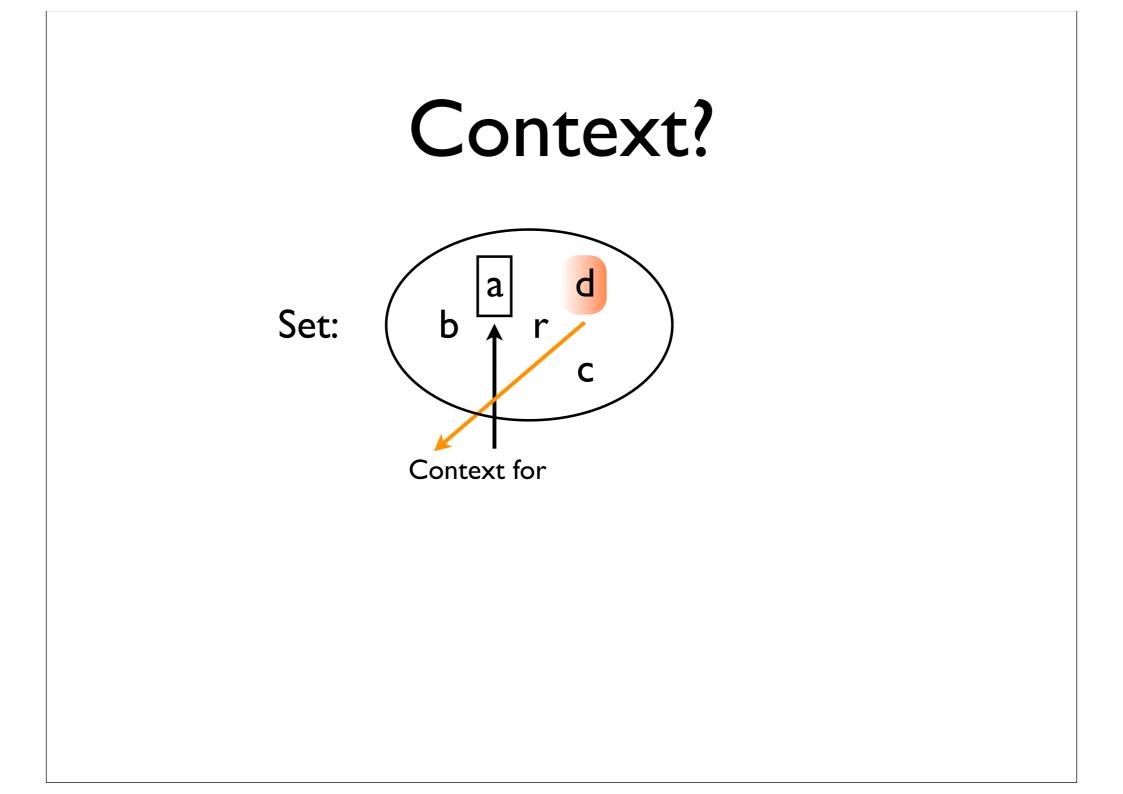


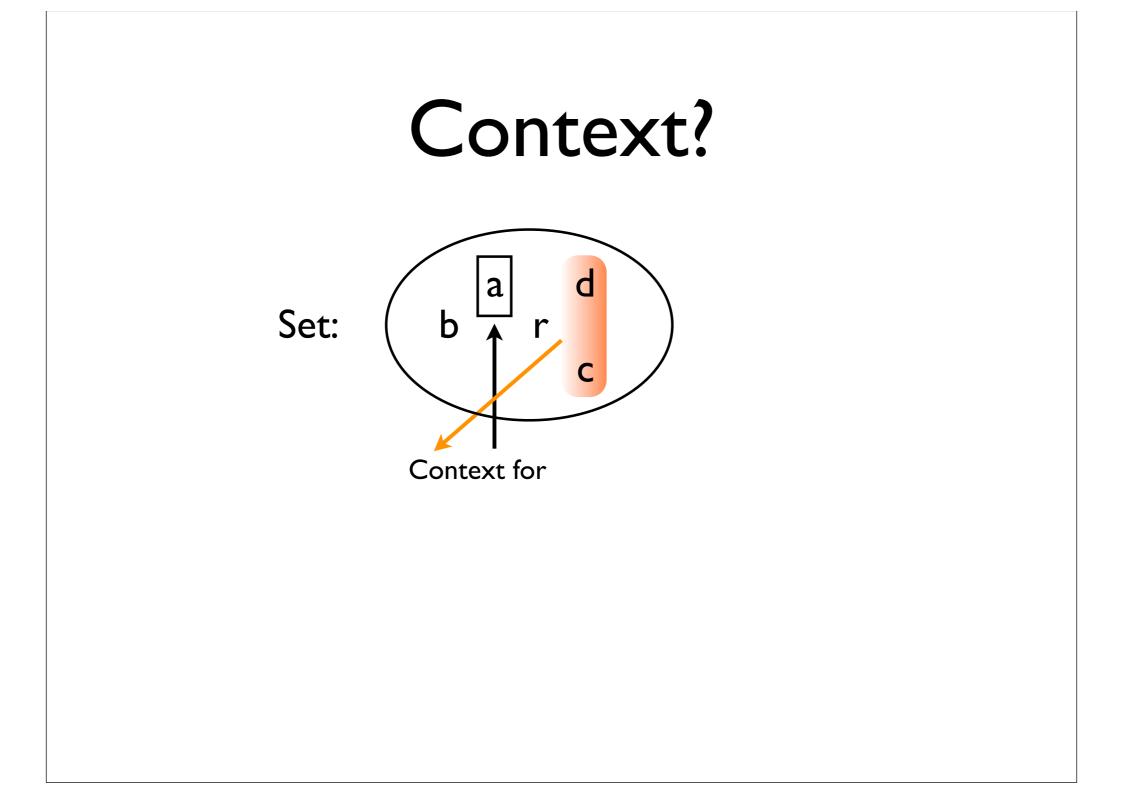


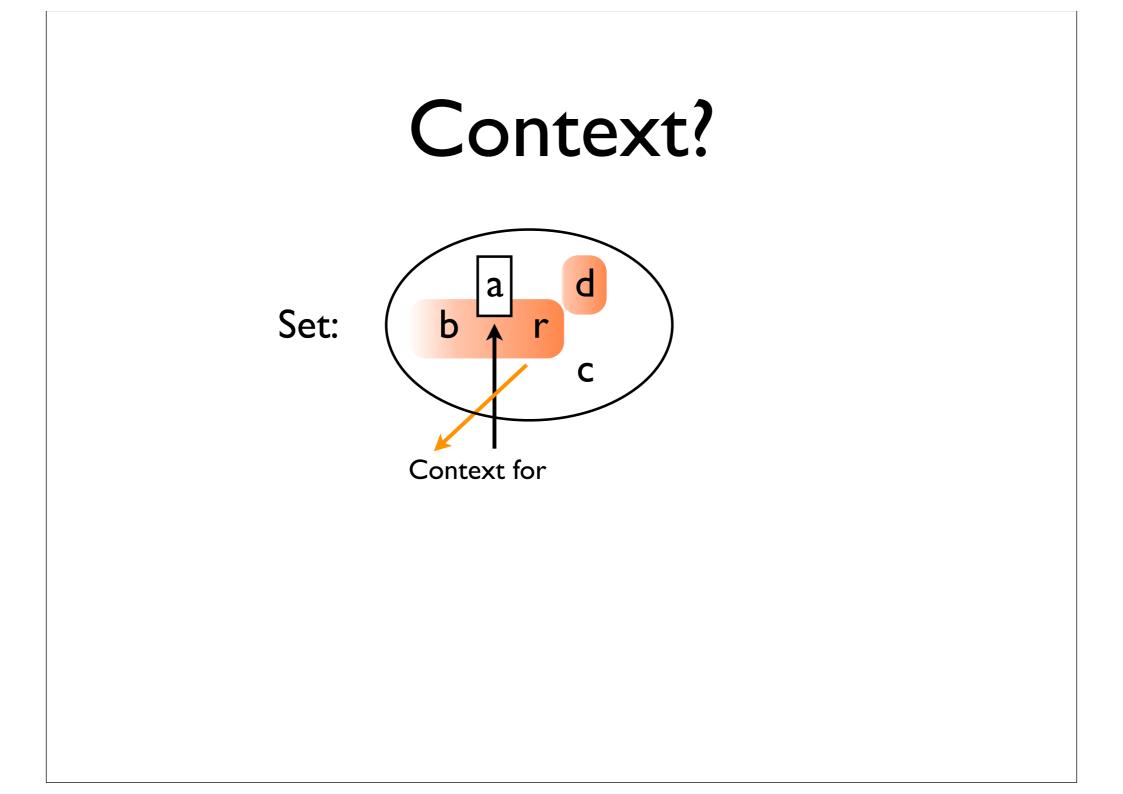










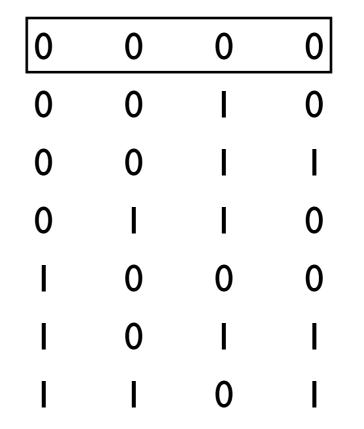


Hinting one path: context partitioning

- Partition into subsets for dependencies:
- ... strong between subsets
- ... weak between elements in same subset
- Condition probabilities on subsets encoded
- Order of subset transmission not important

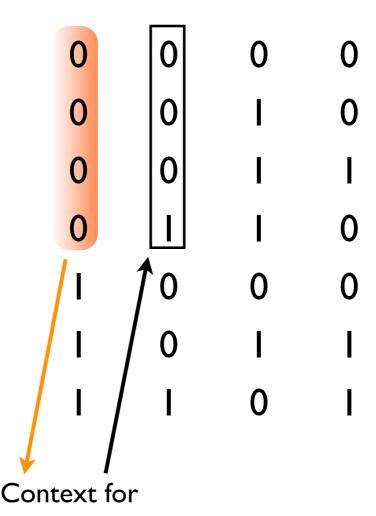
Context in bitwise recursive algorithm

 Represent elements as bitstrings (rows)



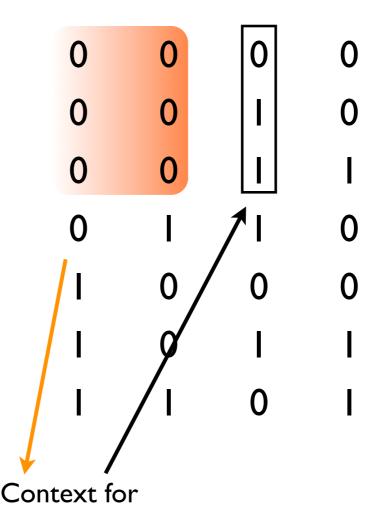
Context in bitwise recursive algorithm

- Represent elements as bitstrings (rows)
- Encode from most to least significant bit
- Context obtained from more significant bits



Context in bitwise recursive algorithm

- Represent elements as bitstrings (rows)
- Encode from most to least significant bit
- Context obtained from more significant bits



Known method 3: Interpolative coding

- Set: { 4, 9, 11, 14, 16, 17, 20, 21 }
- Encode 21 in range [0, |U|)

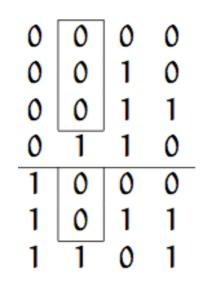
 I4 in range [0, 21)
 9 in range [0, 14)
 4 in range [0, 9)
 I1 in range (9, 14)
 I7 in range (14, 21)
 I6 in range (14, 17) binary
 20 in range (17, 21)

Moffat and Stuiver 2000

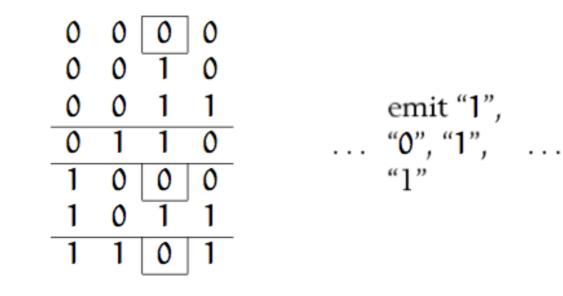
• (Simplified. Can also use known no of elements in range)

Compress set: {1000,0010,0000,1101,0011,1011,0110}

First sort . . .



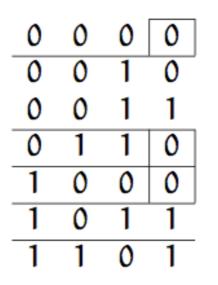
emit "3", "2", ··· recurse ···



emit "1",

"0", "0".

... "1", "1".

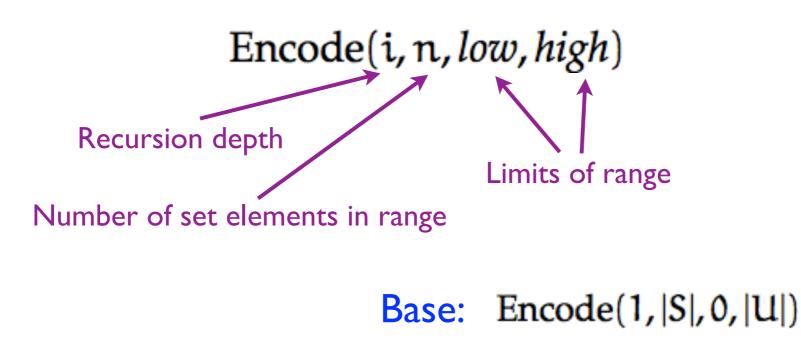


Encode(i, n, low, high):

- If n = 0, no elements remain to encode, and we are done. If n = high low, the elements to be encoded must be low, ..., high - 1, which does not need to be explicitly represented, and again we are done. Otherwise continue:
- 2. Let $b = \lceil \log_2(high low) \rceil$.
- Let m be the number of items among a_i,..., a_{i+n-1} whose bit b − 1 is 0. Since these are the m lower elements of the subarray a_i,..., a_{i+n-1}, m can be found using binary search.
- 4. Output the number m, using some integer encoding (discussed below).
- 5. Recursively invoke $Encode(i, m, low, low + 2^{b-1})$ and $Encode(i + m, n m, low + 2^{b-1}, high)$.

Encoding step

Emit number in smaller range as recursion deepens (~ interpolative)



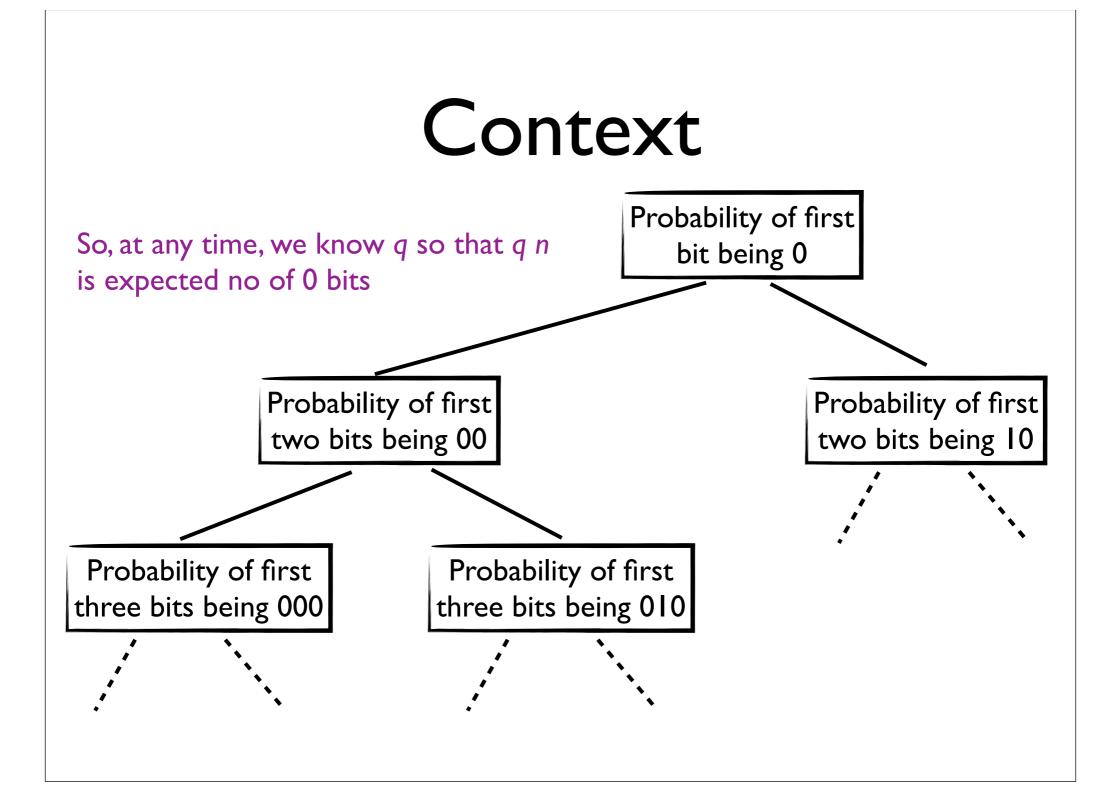
Baseline: uniform element probabilities (no context)

 $s = 2^{\lceil \log_2(high - low) \rceil - 1}$

f = high - low - s

Hypergeometric distribution

 $Pr(elements starting with 0 is m) = \frac{\binom{s}{m}\binom{f}{n-m}}{\binom{s+f}{n}}$



Binomial approximation

Estimate as if draw were with replacement, close if s+f is large in relation to n.

Binomial distribution

Pr(elements starting with 0 is m) = $\binom{n}{m} q^m (1-q)^{n-m}$

Case exclusion: getting rid of nonzero probability for m > s and m < n - f

1. If n > s, reassign, in order, $d \leftarrow n - s$, $n \leftarrow s$, and $f \leftarrow f - d$.

2. Then, if n > f, reassign, in order, $d \leftarrow n - f$, $m \leftarrow m - d$, $n \leftarrow f$, and $s \leftarrow s - d$.

Hypergeometric rescaled

$$\frac{s}{s+f} = q$$

If $s/f \ge q/(1-q)$, reassign $f \leftarrow [s(1-q)/q]$ If s/f < q/(1-q), reassign $s \leftarrow [fq/(1-q)]$

Non-central hypergeometric

- Introduce a weight $w = \frac{f}{s} \cdot \frac{q}{1-q}$
- Wallenius' non-central hypergeometric distribution

		txt/8 orig.	txt/8 order	txt/24 orig.	txt/24 order	words rand.	words order	inverted rand.	inverted order
1	gap	1.71	1.62	2.04	2.03	5.02	4.99	4.64	4.59
2	gap w/o repl.	1.63	1.62	2.04	2.03	5.02	4.99	4.60	4.57
3	interpolative	1.65	1.28	2.16	1.62	5.43	2.20	4.82	2.78
4	dst (Reznik)	2.93	2.86	4.04	3.92	7.29	5.15	6.63	5.36
5	yes/no	1.70	1.70	1.70	1.70	5.09	5.09	5.25	5.25
6	rec. flat	1.99	1.56	2.6 1	2.12	5.60	2.38	5.12	2.95
7	rec. hypergeom.	1.53	1.53	1.96	1.96	5.02	5.02	4.55	4.55
8	rec. binomial	1.23	1.01	1.71	1.46	3.54	2.82	3.26	2.72
9	rec. rescaled hg	1.16	1.01	1.65	1.46	3.48	3.00	3.22	2.81
10	rec. nchg	1.14	1.04	1.62	1.47	N/A	N/A	N/A	N/A
	Sizes:								
11	binary	0.87 (3.00)		0.48 (4.94)		0.02 (14.00)		0.73 (8.00)	
12	uniform	0.92 (3.17)		0.45 (4.64)		0.02 (14.25)		0.77 (8.40)	
13	binomial	0.76 (2.62)		0.34 (3.46)		1.34 (849.86)		1.81 (19.83)	

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