## Encoding and modeling for set compression <br> N. Jesper Larsson (jesl) <br> IT University of Copenhagen (itu. dk)

NII Shonan Meeting, 27-30 Sep. 2013

## Metadata about this work

- Started as broad attack on set compression
- Contribute one algorithm + smaller points
- Not particularly about data structures or time complexity, but setting goals/ideals and finding potentials


## Events, sequences, sets

Event ("circumstance") sequence $E_{1}, \ldots, E_{n}$

$$
\begin{gathered}
\operatorname{Pr}\left(E_{i} \mid E_{1} \cap \cdots \cap E_{i-1}\right) \\
\sum_{i}-\log _{2} \operatorname{Pr}\left(E_{i} \mid E_{1} \cap \cdots \cap E_{i-1}\right) \text { bits }
\end{gathered}
$$

Encoder decoder share premises for what $E_{i}$ mean in terms of specifying message (data)

- Seqence: a certain character is at position $i$
- Set: a certain element is in the set (note: no i)
- Set, alt.: a certain number of elements have a certain property


# Fields of application/ previous work 

- Component (e.g. $\Psi$ )
- Inverted index
- Dictionary
- Data mining (measuring ratio), web graphs, ...


## Universe?

# Encoding a set (or many sets) S, elements drawn from universe $U$ 

## $S \in U$

$|\mathrm{S}|<|\mathrm{U}|$

## Universe?

$$
|\mathrm{S}| \sim \frac{1}{\mathrm{c}}|\mathrm{u}|
$$

small constant

- fixed-length strings?
- characters? patterns?
- probability distribution? (Reznik 201I,
Varshney \& Goyal)

$$
|\mathrm{U}|=\infty
$$



## Narrow focus, for now:

- U may be much larger than S
- Dependencies between elements
- Elements:
- Integers $\in[0,|\mathrm{U}|)$
- = bitstrings of length $\left\lceil\log _{2}|\mathrm{U}|\right\rceil$


## Solved?

- Set: $\{4,9,1 I,|4| 6,, I 7,20,2 I\}$
- Gaps: $\{4,4, \mathrm{I}, 2, \mathrm{I}, 0,2,0\}$
- Geometric distribution
$\operatorname{Pr}($ gap size $k)=(1-p)^{k-1} p, p=\operatorname{Pr}(x \in S)=|S| /|u|$
Optimal code: Golomb (or arithmetic)


## Known (?) method 2: yes/no code

- Arithmetic code for binary source: for each element of $U$, encode whether in $S$
- Estimate $p_{x}=\operatorname{Pr}(x \in S)$
- Use probability ranges $\left[0, p_{x}\right),\left[p_{x}, 1\right)$


## Context?

Sequence: $\underset{\substack{\downarrow \\ \downarrow \\ \text { Context for }}}{b} r$ a c a d a b r a

## Context?



## Context?



## Context?

Sequence: a bracada bracrer

## Context?

Set:


## Context?

Set:


Context for

## Context?

Set:


Context for

## Context?

Set:


Context for

## Context?

Set:


Context for

# Hinting one path: context partitioning 

- Partition into subsets for dependencies:
- ... strong between subsets
- ... weak between elements in same subset
- Condition probabilities on subsets encoded
- Order of subset transmission not important


## Context in bitwise recursive algorithm

- Represent elements as bitstrings (rows)

| 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |

## Context in bitwise recursive algorithm

- Represent elements as bitstrings (rows)
- Encode from most to least significant bit
- Context obtained from more significant bits

| 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |

Context for

## Context in bitwise recursive algorithm

- Represent elements as bitstrings (rows)
- Encode from most to least significant bit
- Context obtained from more significant bits



## Known method 3: Interpolative coding

- Set: $\{4,9, I I, I 4,16, I 7,20,2 I\}$
- Encode 2 I in range $[0,|\mathrm{U}|)$

14 in range $[0,21$ )
9 in range $[0,14)$
4 in range $[0,9)$
II in range $(9,14)$
17 in range $(14,21)$
16 in range $(14,17)$ binary
20 in range $(17,21)$

- (Simplified. Can also use known no of elements in range)


# New method: recursive bitstring set encoding 

Compress set: $\{1000,0010,0000,1101,0011,1011,0110\}$

First sort ...

| 0 | 0 | 0 | 0 |  | 0 | 0 | 0 | 0 |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 0 | 1 | 0 |  | 0 | 0 | 1 | 0 | emit "4", |
| 0 | 0 | 1 | 1 | count 0s | 0 | 0 | 1 | 1 | continue |
| 0 | 1 | 1 | 0 | in first | 0 | 1 | 1 | 0 | recursively |
| 1 | 0 | 0 | 0 | position | 1 | 0 | 0 | 0 | for next |
| 1 | 0 | 1 | 1 |  | 1 | 0 | 1 | 1 | position |
| 1 | 1 | 0 | 1 |  | 1 | 1 | 0 | 1 |  |

# New method: recursive bitstring set encoding 

| 0 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 0 |
| 0 | 0 | 1 | 1 |
| 0 | 0 | 1 | 1 |
| 0 | 1 | 1 | 0 |
| 1 | 0 | 0 | 0 |
| 1 | 0 | 1 | 1 |
| 1 | 1 | 0 | 1 |

emit " 3 ", " 2 ", $\ldots$
recurse

## New method: recursive bitstring set encoding



## New method: recursive bitstring set encoding

| 0 | 0 | 0 | 0 |  |
| :--- | :--- | :--- | :--- | :--- |
| 0 | 0 | 1 | 0 |  |
| 0 | 0 | 1 | 1 |  |
| 0 | 1 | 1 | 0 |  |
| 1 | 0 | 0 | 0 |  |
| 1 | 0 | 1 | 1 | emit " 11 ", "1", " 1 ", |
| 1 | 1 | 0 | 1 | $" 0$ ", " 0 ", |

## New method: recursive bitstring set encoding

Encode(i, n, low, high):

1. If $\mathrm{n}=0$, no elements remain to encode, and we are done. If $\mathrm{n}=$ high - low, the elements to be encoded must be low, ..., high -1 , which does not need to be explicitly represented, and again we are done. Otherwise continue:
2. Let $\mathrm{b}=\left\lceil\log _{2}(\right.$ high $\left.-l o w)\right\rceil$.
3. Let $m$ be the number of items among $a_{i}, \ldots, a_{i+n-1}$ whose bit $b-1$ is 0 . Since these are the $m$ lower elements of the subarray $a_{i}, \ldots, a_{i+n-1}, m$ can be found using binary search.
4. Output the number $m$, using some integer encoding (discussed below).
5. Recursively invoke Encode $\left(i, m\right.$, low, low $\left.+2^{b-1}\right)$ and Encode $(i+m, n-m, l o w+$ $2^{\mathrm{b}-1}$,high).

## Encoding step

Emit number in smaller range as recursion deepens ( $\sim$ interpolative)


Number of set elements in range

Base: Encode(1,|S|,0,|U|)

# Baseline: uniform element probabilities (no context) 

$$
\begin{gathered}
s=2^{\left[\log _{2}(h i g h-l o w)\right\rceil-1} \\
\mathrm{f}=h i g h-l o w-\mathrm{s}
\end{gathered}
$$

Hypergeometric distribution
$\operatorname{Pr}($ elements starting with 0 is $m)=\frac{\binom{s}{m}\binom{f}{n-m}}{\binom{s+f}{n}}$

## Context



## Binomial approximation

Estimate as if draw were with replacement, close if $s+f$ is large in relation to $n$.

## Binomial distribution

$\operatorname{Pr}($ elements starting with 0 is $\mathfrak{m})=\binom{n}{m} q^{m}(1-q)^{n-m}$

Case exclusion: getting rid of nonzero probability for $m>s$ and $m<n-f$

1. If $n>s$, reassign, in order, $d \leftarrow n-s, n \leftarrow s$, and $f \leftarrow f-d$.
2. Then, if $n>f$, reassign, in order, $d \leftarrow n-f, m \leftarrow m-d, n \leftarrow f$, and $s \leftarrow s-d$.

## Hypergeometric rescaled

$$
\frac{s}{s+f}=q
$$

If $s / f \geq q /(1-q)$, reassign $f \leftarrow[s(1-q) / q]$
If $s / f<q /(1-q)$, reassign $s \leftarrow[f q /(1-q)]$

## Non-central hypergeometric

- Introduce a weight $w=\frac{f}{s} \cdot \frac{q}{1-q}$
- Wallenius' non-central hypergeometric distribution


## Results

|  | $t x t / 8 \quad t x t / 8$ orig. order | $\begin{gathered} \text { txt/24 } \\ \text { orig. } \end{gathered}$ | $\begin{gathered} \text { txt/24 } \\ \text { order } \end{gathered}$ | words <br> rand. | words order | inverted rand. | inverted order |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| gap | 1.711 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.64 | 4.59 |
| 2 gap w/o repl. | 1.631 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.60 | 4.57 |
| 3 interpolative | 1.651 .28 | 2.16 | 1.62 | 5.43 | 2.20 | 4.82 | 2.78 |
| 4 dst (Reznik) | 2.932 .86 | 4.04 | 3.92 | 7.29 | 5.15 | 6.63 | 5.36 |
| 5 yes/no | 1.701 .70 | 1.70 | 1.70 | 5.09 | 5.09 | 5.25 | 5.25 |
| 6 rec. flat | 1.991 .56 | 2.61 | 2.12 | 5.60 | 2.38 | 5.12 | 2.95 |
| 7 rec. hypergeom. | 1.531 .53 | 1.96 | 1.96 | 5.02 | 5.02 | 4.55 | 4.55 |
| 8 rec. binomial | 1.231 .01 | 1.71 | 1.46 | 3.54 | 2.82 | 3.26 | 2.72 |
| 9 rec. rescaled hg | 1.161 .01 | 1.65 | 1.46 | 3.48 | 3.00 | 3.22 | 2.81 |
| 10 rec. nchg | 1.141 .04 | 1.62 | 1.47 | N/A | N/A | N/A | N/A |
| Sizes: |  |  |  |  |  | 0.73 (8.00) |  |
| 11 binary | 0.87 (3.00) | 0.48 (4.94) |  | 0.02 (14.00) |  |  |  |
| 12 uniform | 0.92 (3.17) | 0.45 (4.64) |  | 0.02 (14.25) |  | 0.77 (8.40) |  |
| 13 binomial | 0.76 (2.62) | 0.34 | 3.46) | 1.34 (8) | 49.86) | 1.81 (19.83) |  |

## Results

|  | $t x t / 8 \quad t x t / 8$ orig. order | txt/24 orig. | txt/24 order | words <br> rand. | words order | inverted rand. | inverted order |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 gap | 1.711 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.64 | 4.59 |
| 2 gap w/o repl. | 1.631 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.60 | 4.57 |
| 3 interpolative | 1.651 .28 | 2.16 | 1.62 | 5.43 | 2.20 | 4.82 | 2.78 |
| 4 dst (Reznik) | 2.932 .86 | 4.04 | 3.92 | 7.29 | 5.15 | 6.63 | 5.36 |
| 5 yes/no | 1.701 .70 | 1.70 | 1.70 | 5.09 | 5.09 | 5.25 | 5.25 |
| 6 rec. flat | 1.991 .56 | 2.61 | 2.12 | 5.60 | 2.38 | 5.12 | 2.95 |
| 7 rec. hypergeom. | 1.531 .53 | 1.96 | 1.96 | 5.02 | 5.02 | 4.55 | 4.55 |
| 8 rec. binomial | 1.231 .01 | 1.71 | 1.46 | 3.54 | 2.82 | 3.26 | 2.72 |
| 9 rec. rescaled hg | 1.161 .01 | 1.65 | 1.46 | 3.48 | 3.00 | 3.22 | 2.81 |
| 10 rec. nchg | 1.141 .04 | 1.62 | 1.47 | N/A | N/A | N/A | N/A |
| Sizes: |  | 0.48 (4.94) |  | 0.02 (14.00) |  | 0.73 (8.00) |  |
| 11 binary | 0.87 (3.00) |  |  |  |  |  |  |
| 12 uniform | 0.92 (3.17) | 0.45 (4.64) |  | 0.02 (14.25) |  | 0.77 (8.40) |  |
| 13 binomial | 0.76 (2.62) | 0.34 | 3.46) | 1.34 (8) | 49.86) | 1.81 (19.83) |  |

## Results

|  | txt/8 txt/8 orig. order | $\begin{gathered} \text { txt/24 } \\ \text { orig. } \end{gathered}$ | $\begin{aligned} & \text { txt/24 } \\ & \text { order } \end{aligned}$ | words rand. | words order | inverted rand. | inverted order |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| gap | 1.711 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.64 | 4.59 |
| 2 gap w/o repl. | 1.631 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.60 | 4.57 |
| 3 interpolative | 1.651 .28  <br> 2.  | 2.16 | 1.62 | 5.43 | 2.20 | 4.82 | 2.78 |
| 4 dst (Reznik) | $2.93 \quad 2.86$ | 4.04 | 3.92 | 7.29 | 5.15 | 6.63 | 5.36 |
| 5 yes/no | 1.701 .70 | 1.70 | 1.70 | 5.09 | 5.09 | 5.25 | 5.25 |
| 6 rec. flat | 1.991 .56 | 2.61 | 2.12 | 5.60 | 2.38 | 5.12 | 2.95 |
| 7 rec. hypergeom. | 1.531 .53 | 1.96 | 1.96 | 5.02 | 5.02 | 4.55 | 4.55 |
| 8 rec. binomial | 1.231 .01 | 1.71 | 1.46 | 3.54 | 2.82 | 3.26 | 2.72 |
| 9 rec. rescaled hg | 1.161 .01 | 1.65 | 1.46 | 3.48 | 3.00 | 3.22 | 2.81 |
| 10 rec. nchg | 1.141 .04 | 1.62 | 1.47 | N/A | N/A | N/A | N/A |
| Sizes: |  |  |  |  |  |  |  |
| 11 binary | 0.87 (3.00) | 0.48 | (4.94) | 0.02 | 14.00) | 0.73 | 8.00) |
| 12 uniform | 0.92 (3.17) | 0.45 | (4.64) | 0.02 | (14.25) | 0.77 | 8.40) |
| 13 binomial | 0.76 (2.62) | 0.34 | (3.46) | 1.34 (8) | (89.86) | 1.81 | (19.83) |

## Results



## Results

|  | $t x t / 8 \quad t x t / 8$ orig. order | $\begin{gathered} \text { txt/24 } \\ \text { orig. } \end{gathered}$ | $\begin{gathered} \text { txt/24 } \\ \text { order } \end{gathered}$ | words <br> rand. | words order | inverted rand. | inverted order |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| gap | 1.711 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.64 | 4.59 |
| 2 gap w/o repl. | 1.631 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.60 | 4.57 |
| 3 interpolative | 1.651 .28 | 2.16 | 1.62 | 5.43 | 2.20 | 4.82 | 2.78 |
| 4 dst (Reznik) | 2.932 .86 | 4.04 | 3.92 | 7.29 | 5.15 | 6.63 | 5.36 |
| 5 yes/no | 1.701 .70 | 1.70 | 1.70 | 5.09 | 5.09 | 5.25 | 5.25 |
| 6 rec. flat | 1.991 .56 | 2.61 | 2.12 | 5.60 | 2.38 | 5.12 | 2.95 |
| 7 rec. hypergeom. | 1.531 .53 | 1.96 | 1.96 | 5.02 | 5.02 | 4.55 | 4.55 |
| 8 rec. binomial | 1.231 .01 | 1.71 | 1.46 | 3.54 | 2.82 | 3.26 | 2.72 |
| 9 rec. rescaled hg | 1.161 .01 | 1.65 | 1.46 | 3.48 | 3.00 | 3.22 | 2.81 |
| 10 rec. nchg | 1.141 .04 | 1.62 | 1.47 | N/A | N/A | N/A | N/A |
| Sizes: |  |  |  |  |  | 0.73 (8.00) |  |
| 11 binary | 0.87 (3.00) | 0.48 (4.94) |  | 0.02 (14.00) |  |  |  |
| 12 uniform | 0.92 (3.17) | 0.45 (4.64) |  | 0.02 (14.25) |  | 0.77 (8.40) |  |
| 13 binomial | 0.76 (2.62) | 0.34 | 3.46) | 1.34 (8) | 49.86) | 1.81 (19.83) |  |

## Results

|  | $t x t / 8 \quad t x t / 8$ orig. order | $\begin{gathered} \text { txt/24 } \\ \text { orig. } \end{gathered}$ | $\begin{aligned} & \text { txt } / 24 \\ & \text { order } \end{aligned}$ | words rand. | words order | inverted rand. | inverted order |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| gap | 1.711 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.64 | 4.59 |
| gap w/o repl. | 1.631 .62 | 2.04 | 2.03 | 5.02 | 4.99 | 4.60 | 4.57 |
| interpolative | 1.65 1.28 <br> 1.68  | 2.16 | 1.62 | 5.43 | 2.20 | 4.82 | 2.78 |
| 4 dst (Reznik) | $2.93 \quad 2.86$ | 4.04 | 3.92 | 7.29 | 5.15 | 6.63 | 5.36 |
| 5 yes/no | $1.70 \quad 1.70$ | 1.70 | 1.70 | 5.09 | 5.09 | 5.25 | 5.25 |
| 6 rec. flat | 1.991 .56 | 2.61 | 2.12 | 5.60 | 2.38 | 5.12 | 2.95 |
| rec. hypergeom. | 1.53 1.53 | 1.96 | 1.96 | 5.02 | 5.02 | 4.55 | 4.55 |
| 8 rec. binomial | 1.231 .01 | 1.71 | 1.46 | 3.54 | 2.82 | 3.26 | 2.72 |
| 9 rec. rescaled hg | 1.161 .01 | 1.65 | 1.46 | 3.48 | 3.00 | 3.22 | 2.81 |
| 10 rec. nchg | 1.141 .04 | 1.62 | 1.47 | N/A | N/A | N/A | N/A |
| Sizes: |  |  |  |  |  |  |  |
| 11 binary | 0.87 (3.00) | 0.48 | 4.94) | 0.02 | (14.00) | 0.73 | 8.00) |
| 12 uniform | 0.92 (3.17) | 0.45 | (4.64) | 0.02 | (14.25) | 0.77 | 8.40) |
| 13 binomial | 0.76 (2.62) | 0.34 | (3.46) | 1.34 | 849.86) | 1.81 | (19.83) |

## Thank you.

